



Analysis and Design of Algorithms

UNIT-1

Recurrence Relations

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Conclusion

Analysis of Algorithm

Algorithm Analysis → It provides background information that gives us a general idea of how long an Algo will take for a given problem set. Suppose a problem have n inputs. If we get comparison n times for arrange them into ascending order so $n \times n$ operations performed.

- purpose of Algo is "not to find a formula".
- Analysis of Algo without regard to any specific computer type.

Algo → Counts the No. of Different characters in a file.

For all 256 char do

Assign zero to the counter] loop 1
end for loop.

While there are more characters in the file do

Get the next character

Increment the counter for this character by one

end while loop.

] loop 2

Description → - total 256 passes for the initialization loop is loop 1

- If n char in the input file there are n passes for the second loop ~~is~~ loop 2.

- For each pass of the loop there is a check exist that x for loop variable is ~~within~~ within the bounds.

- Initialization loop (i.e. loop 1) does a set of 257 assignments

- 256 increments for the loop variable

- 257 checks that this variable is within the loop bounds.

For the second loop, we will need to do check of the condition $N+1$ times (+1 for the last check when the file is empty.) & we will increment N counters.

The total No. of operations is

- Increment $N+256$
- Assignments 257
- checks $(N+258)$ conditions.

So ~~if~~ increase char in file ie \rightarrow 500, 1000, 10000 etc
we have to increase the operations & checks.

So Algo. Analysis requires a set of Rules to determine how operations are to be counted.

Exact Analysis Rules \rightarrow

- (1) Assume A time unit.
- (2) Execution of one of the following operations take time 1 (one):
 - (a) Assignment operations
 - (b) I/O operation (single)
 - (c) Single Boolean operation, numeric operation
 - (d) Single Arithmetic operations
 - (e) function Return
 - (f) Array Index operations, pointer Dereference

(3) Running time of a selection statement (switch) is the time for the condition evaluation + Maximum running time for the individual clauses in the selection.

④ Loop execution time is the sum of the body loop over a number the is executed + time for the loop check + update operations + time for the loop return.

⑤ Running time of a function call is 1 for setup + The time for any parameter calculations + Time Required for the execution of the function body.

The analysis of an Algo. is to evaluate the performance of the Algo. based on the given models + metrics.

① Input size.

② Running time \rightarrow worst case + Average case
generally we finding only the worst case Running time.

③ Order of growth \rightarrow we generally see the growth rate of the Running time.

- We only consider the leading terms of a time formula.

The leading term is n^2 in expression $n^2 + 100n + 5000$.

So Algo. Analysis is a part of Computational Complexity Theory. In Theoretical Analysis of Algo. It is common to estimate their complexity in "Asymptotic

sense" ie \rightarrow To estimate the complexity function for large length of input.

\rightarrow Usually Asymptotic estimates are used because different implementation of the same Algo. may differ in efficiency.

Kinds of analyses

Worst-case: (usually)

- $T(n)$ = maximum time of algorithm on any input of size n .

Average-case: (sometimes)

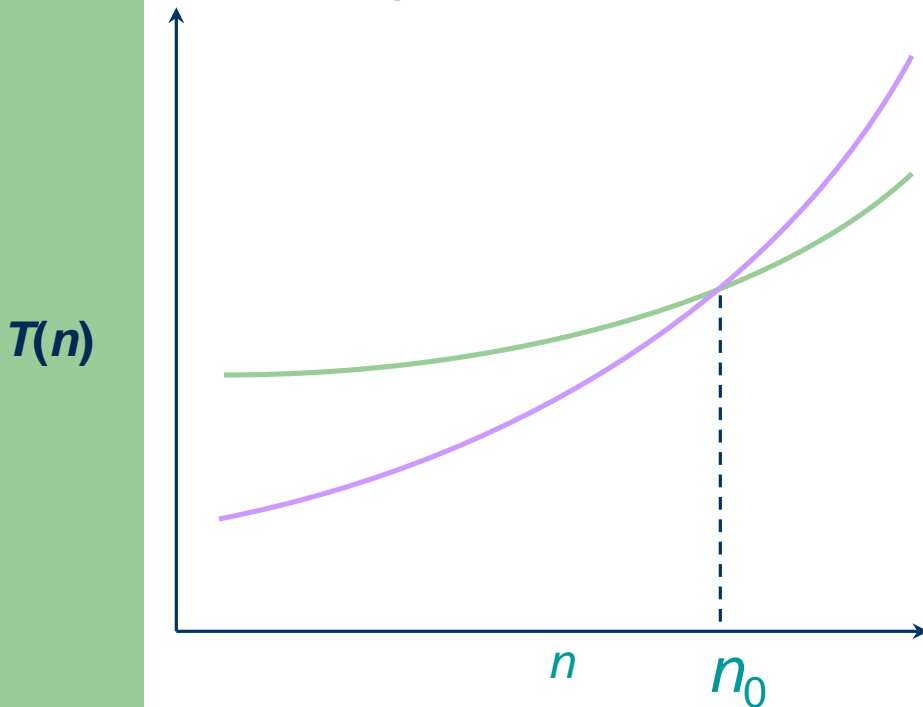
- $T(n)$ = expected time of algorithm over all inputs of size n .
- Need assumption of statistical distribution of inputs.

Best-case: (NEVER)

- Cheat with a slow algorithm that works fast on *some* input.

Asymptotic performance

When n gets large enough, a $\Theta(n^2)$ algorithm *always* beats a $\Theta(n^3)$ algorithm.



- Asymptotic analysis is a useful tool to help to structure our thinking toward better algorithm
- We shouldn't ignore asymptotically slower algorithms, however.
- Real-world design situations often call for a careful balancing



Recurrence Relation

Sequences and Recurrence Relations

Consider the following two sequences:

$$S_1 : 3, 5, 7, 9, \dots$$

$$S_2 : 3, 9, 27, 81, \dots$$

We can find a formula for the n th term of sequences S_1 and S_2 by observing the pattern of the sequences.

$$S_1 : 2 \cdot 1 + 1, 2 \cdot 2 + 1, 2 \cdot 3 + 1, 2 \cdot 4 + 1, \dots$$

$$S_2 : 3^1, 3^2, 3^3, 3^4, \dots$$

For S_1 , $a_n = 2n + 1$ for $n \geq 1$, and for S_2 , $a_n = 3^n$ for $n \geq 1$. This type of formula is called an **explicit formula** for the sequence, because using this formula we can directly find any term of the sequence without using other terms of the sequence. For example, $a_3 = 2 \cdot 3 + 1 = 7$.

Sequences and Recurrence Relations

A **recurrence relation** for a sequence $a_0, a_1, a_2, \dots, a_n, \dots$ is an equation that relates a_n to some of the terms $a_0, a_1, a_2, \dots, a_{n-2}, a_{n-1}$ for all integers n with $n \geq k$, where k is a nonnegative integer. The **initial conditions** for the recurrence relation are a set of values that explicitly define some of the members of $a_0, a_1, a_2, \dots, a_{k-1}$.

The equation

$$a_n = 2a_{n-1} + a_{n-2} \quad \text{for all } n \geq 2,$$

as defined above, relates a_n to a_{n-1} and a_{n-2} . Here $k = 2$. So this is a recurrence relation with initial conditions $a_0 = 5$ and $a_1 = 7$.

Recursion and Recurrences

- **Recursion** is a particularly powerful kind of **reduction**, which can be described loosely as follows:
 - If the given instance of the problem is small or simple enough, just solve it.
 - Otherwise, reduce the problem to one or more simpler instances of the same problem.
- **Recursion** is generally **expressed in terms of recurrences**.
- In other words, when an **algorithm calls to itself**, we can often **describe its running time by a recurrence equation**.
- **recurrence equation** describes the **overall running time of a problem of size n in terms of the running time on smaller inputs**.

Recursive Algorithms

- A recursive algorithm is one in which objects are defined in terms of other objects of the same type.
- Advantages:
 - Simplicity of code
 - Easy to understand
- Disadvantages
 - Memory
 - Speed
 - Possibly redundant work
- Tail recursion offers a solution to the memory problem, but really, do we need recursion?

Recursive Algorithms: Analysis

- We have already discussed how to analyze the running time of (iterative) algorithms
- To analyze recursive algorithms, we require more sophisticated techniques
- Specifically, we study how to defined & solve recurrence relations

Motivating Examples: Factorial

- Recall the factorial function:

$$n! = \begin{cases} 1 & \text{if } n = 1 \\ n \cdot (n-1) & \text{if } n > 1 \end{cases}$$

- Consider the following (recursive) algorithm for computing $n!$

FACTORIAL

Input: $n \in \mathbb{N}$

Output: $n!$

- If** $(n=1)$ or $(n=0)$
- Then Return** 1
- Else Return** $n \times \text{FACTORIAL}(n-1)$
- Endif**
- End**

Factorial: Analysis

How many multiplications $M(x)$ does factorial perform?

- When $n=1$ we don't perform any
- Otherwise, we perform one...
- ... plus how ever many multiplications we perform in the recursive call `FACTORIAL(n-1)`
- The number of multiplications can be expressed as a formula (similar to the definition of $n!$)

$$M(0) = 0$$

$$M(n) = 1 + M(n-1)$$

- This relation is known as a recurrence relation

Recurrence Relation

- A **recurrence relation** for the sequence, a_0, a_1, \dots, a_n , is an equation that relates a_n to certain of its predecessors a_0, a_1, \dots, a_{n-1} .
- Initial conditions for the sequence a_0, a_1, \dots are explicitly given values for a finite number of the terms of the sequence.

Recurrence Relations

- **Definition:** A recurrence relation for a sequence $\{a_n\}$ is an equation that expresses a_n in terms of one or more of the previous terms in the sequence:

$$a_0, a_1, a_2, \dots, a_{n-1}$$

for all integers $n \geq n_0$ where n_0 is a non-negative integer.

- A sequence is called a solution of a recurrence if its terms satisfy the recurrence relation

Recurrence Relations: Solutions

- Consider the recurrence relation -

$$a_n = 2 * a_{n-1} - a_{n-2}$$

- It has the following sequences a_n as **solutions**
 - $a_n = 3n$
 - $a_n = n+1$
 - $a_n = 5$
- The **initial conditions + recurrence relation** uniquely determine the **sequence**.