



Analysis and Design of Algorithms

UNIT-1

Recurrence Relations

Content

- Solving Recurrence Relations
 - Substitution Method



SUBSTITUTION METHOD

2. SUBSTITUTION METHOD

- The substitution method comprises of 3 steps-
 1. “Making a good guess” method or **Guess the form of the solution.**
 2. then use induction or **Verify by induction method.**
 3. Solve for constants or **to find the constants and show that solution works.**
- Examples:
 - $T(n) = 2T(n/2) + \Theta(n) \rightarrow T(n) = \Theta(n \log n)$

2. SUBSTITUTION METHOD

- *The substitution method can be used to establish either upper or lower bounds on a recurrence.*
- We **substitute the guessed solution** for the function, **when applying the inductive hypothesis** to smaller values. Hence the name “substitution method”.
- This method is powerful, but **we must be able to guess the form of the answer in order to apply it.**
 - ▣ In this method – “To guess the solution, **play around with small values for insight**”.
 - ▣ That is in this method first one start “guess a solution and prove by induction”.

Substitution method - An example

- $T(n) = 2T(n/2) + n$

We guess that the solution is $T(n) = O(n \lg n)$.

i.e. to show that $T(n) \leq c n \lg n$, for some constant $c > 0$ and $n \geq m$.

Assume that this bound holds for $[n/2]$. So, we get

$$T(n) \leq 2(c * (n/2) \lg((n/2))) + n \quad \rightarrow T(n/2) = c * (n/2) \lg((n/2))$$

$$\leq cn \lg(n/2) + n$$

$$= cn \lg n - cn \lg 2 + n$$

$$= cn \lg n - cn + n$$

$$\leq cn \lg n$$

where, the last step holds as long as $c \geq 1$.