

Breadth first Search :- Unweighted Graph

①

Input: unweighted graph and start vertex u .

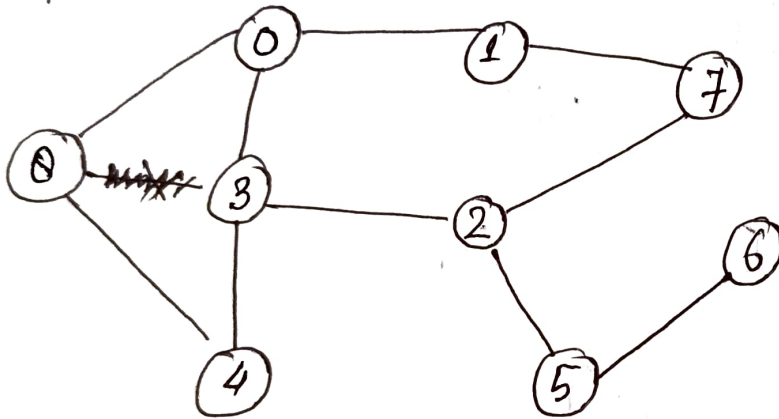
Idea: maintain a set R of vertices that have been reached but not searched and a set S of vertices that have been searched.

set $R \{ \}$ $\rightarrow$ use FIFO (queue)

Initialization: $R = \{u\}$, $S = \phi$, $d(u, u) = 0$

Iteration: $R \neq \phi$, search from the vertex v of R . The neighbors of v not in $S \cup R$ are added to R & assigned distance $d(u, v) + 1$, and then v is removed from R & placed in S .

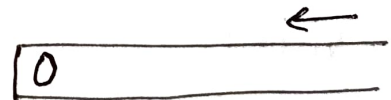
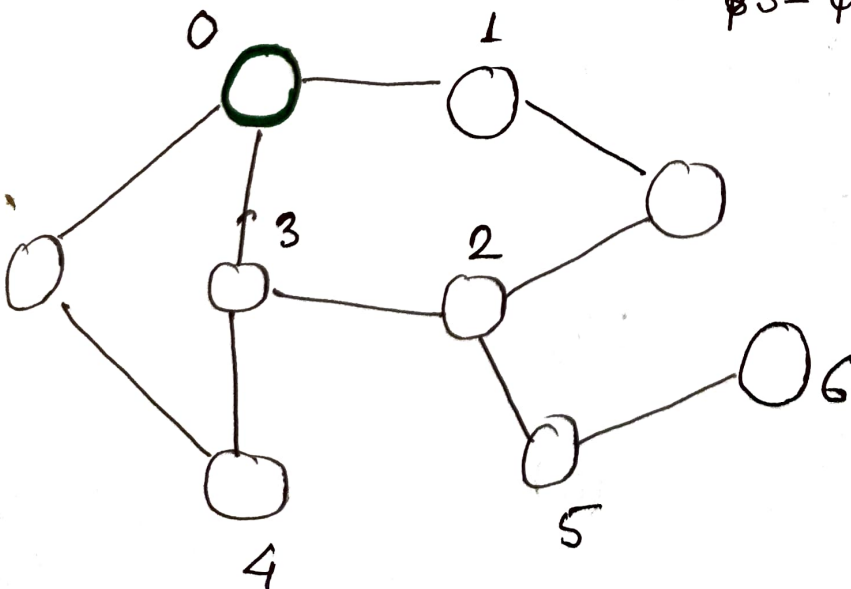
Example $\rightarrow$



* green outlined ~~circle~~ refers queue

* Red internal refers visited.

Step 1:- $R = \{0\}$, 0 is starting vertex. Now maintain a queue
 $S = \phi$, $d(0, 0) = 0$



0 is enqueue.

Step 2: search neighbors of 0 vertex & those are placed in queue
0 is visited

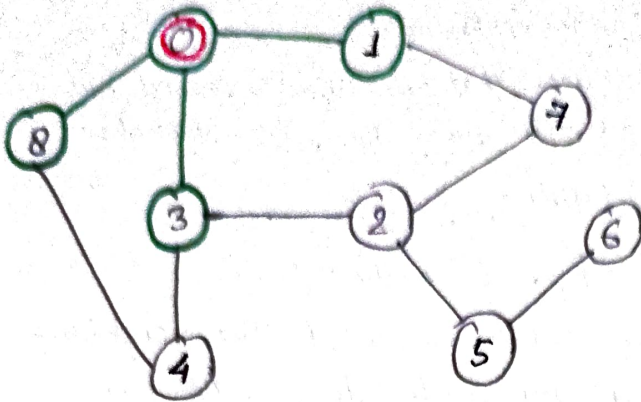
$$R = \{0\}$$

$$S = \{1, 3, 8\}$$

$$d(0,0) = 0$$

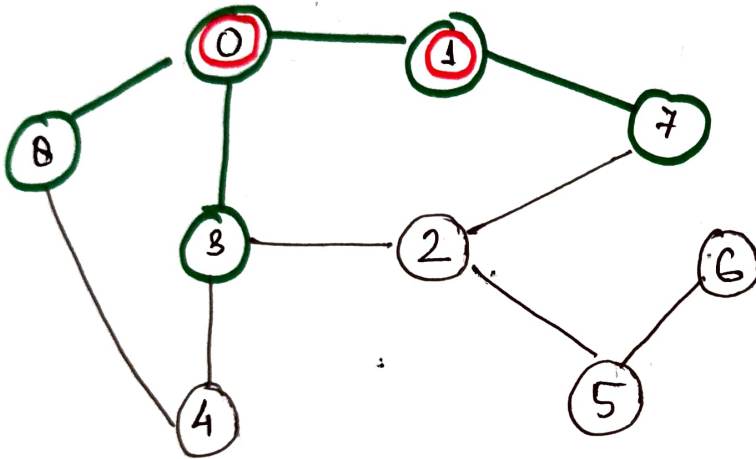
1 3 8

Now next node is 1



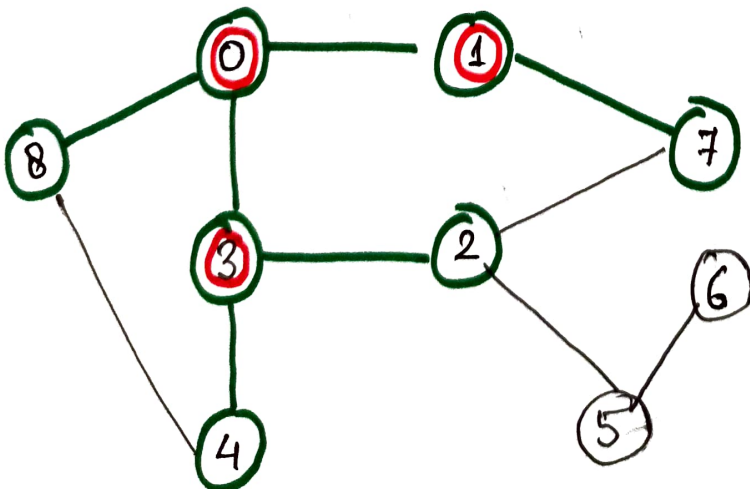
Step 3:- $R = \{0, 1\}$, $S = \{3, 8, 7\}$, $d(0,1) = 1$,
1 is visited

3 8 7



Step 4:- $R = \{0, 1, 3\}$, $S = \{8, 7, 2, 4\}$, $d(0,3) = 1$,
3 is visited

8 7 2 4

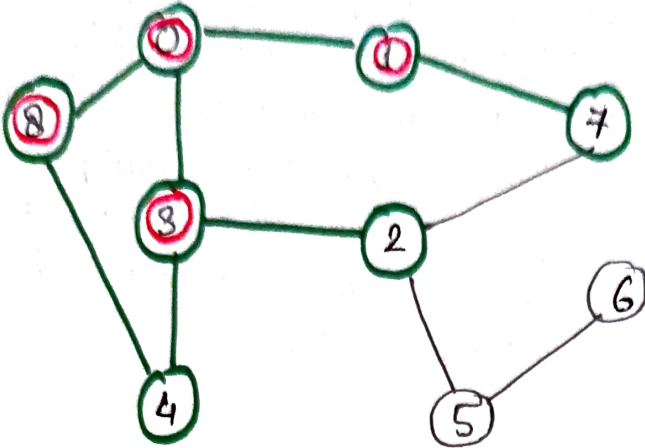


(3)

Step 5 $\rightarrow R = \{0, 1, 3, 8\}$, $S = \{7, 2, 4\}$, $d(0, 8) = 1$

8 is visited

1 7 2 4



Step 6 $\rightarrow R = \{0, 1, 3, 8, 7\}$,

$S = \{2, 4\}$, $d(0, 7) = 2$, 7 is visited

Step 7 $\rightarrow R = \{0, 1, 3, 8, 7, 2\}$

$S = \{4, 5\}$, $d(0, 2) = 2$, 2 visited

Step 8 $\rightarrow R = \{0, 1, 3, 8, 7, 2, 4\}$

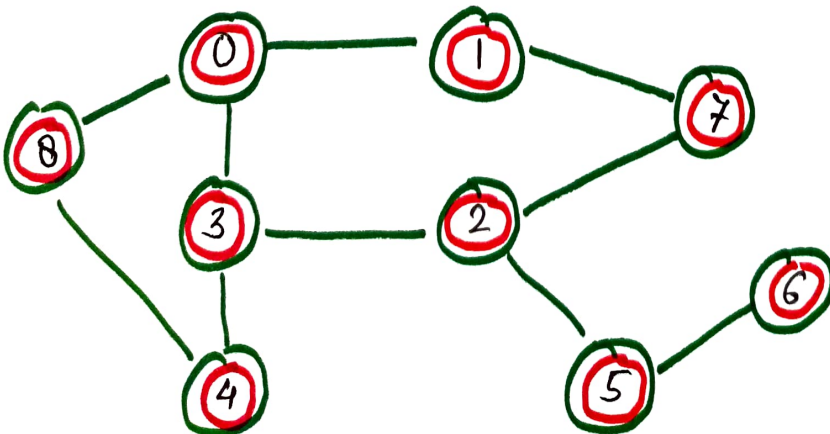
$S = \{5\}$, $d(0, 4) = 2$, 4 visited

Step 9 $\rightarrow R = \{0, 1, 3, 8, 7, 2, 4, 5\}$

$S = \{6\}$, $d(0, 5) = 3$, 5 visited

Step 10 $\rightarrow R = \{0, 1, 3, 8, 7, 2, 4, 5, 6\}$

$S = \{\}$, $d(0, 6) = 4$, 6 visited.



Dijkstra's Algorithm — for weighted graph {only pos. weights}

Input: — A graph with non-negative edge weights & starting vertex u . The weight of edge xy is $w(x, y)$; let $w(x, y) = \infty$, if xy is not an edge (or not participating)

Idea → Maintain the set S of vertices to which the shortest path from u is known, enlarging S to include all vertices. To do this, maintain tentative distance $t(z)$ from u to each $z \notin S$, being the length of the shortest u, z -path.

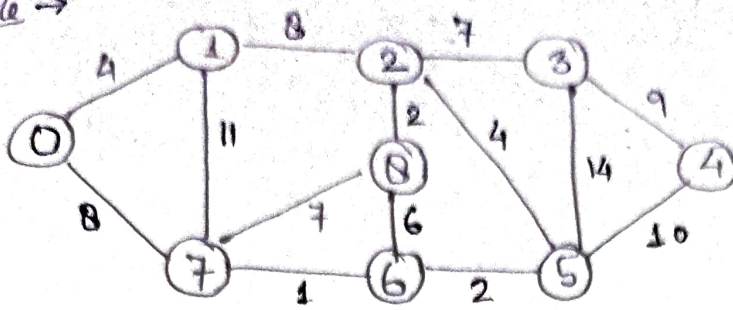
Steps →

- ① Mark ~~self~~ selected initial node with a current distance of 0 & rest with infinity.
- ② Set the non-visited node with smallest current distance as current node C .
- ③ for each neighbor N of current node C : add the current distance of C with weight of the edge connecting $C-N$. If it's smaller than the current distance of N , set it as the new current distance of N .
- ④ Mark the current node C as visited.
- ⑤ Repeat the process from step 2.

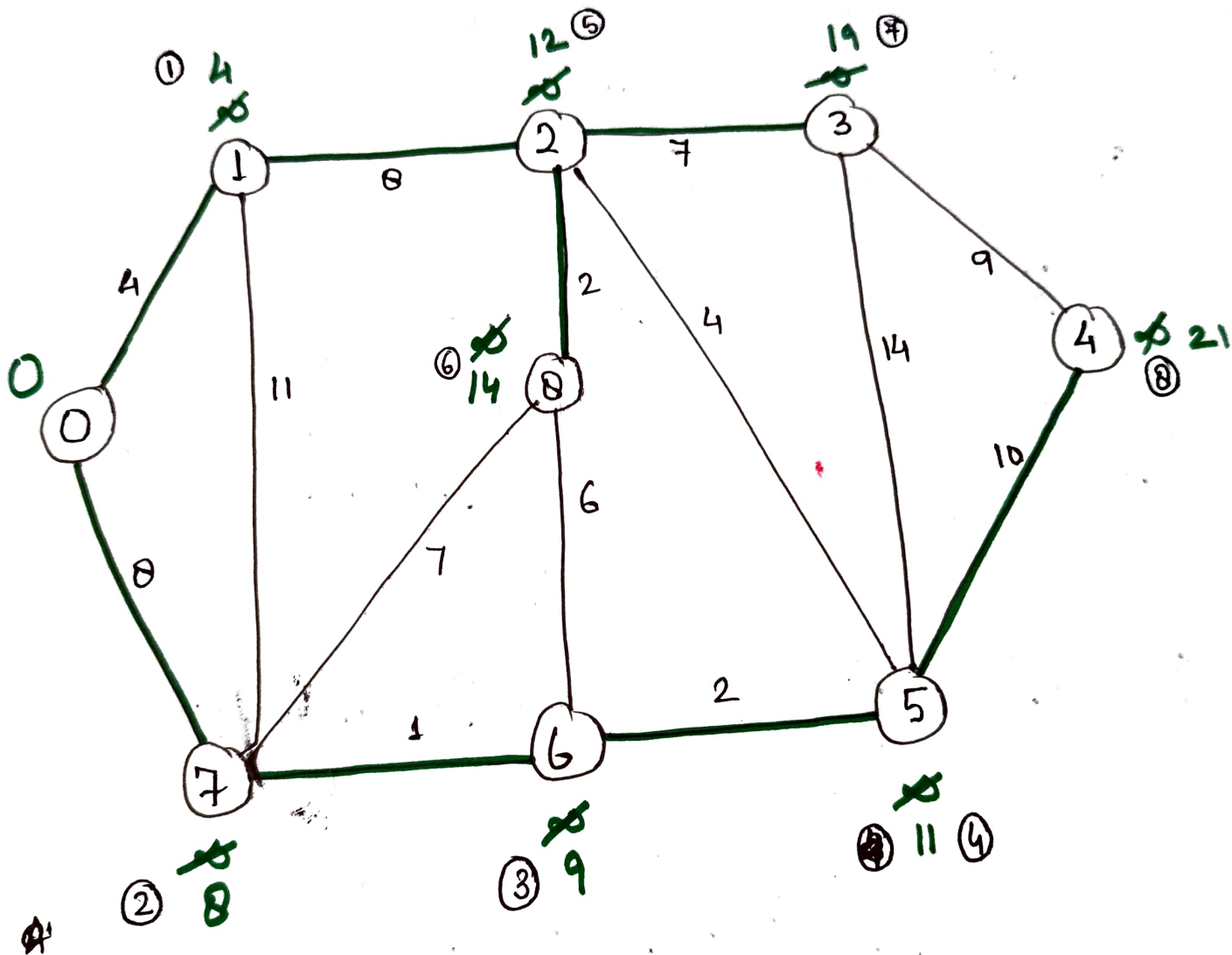
In Dijkstra's Algo we are using priority queue.

~~$i = \text{EXTRACTMIN}; \text{mark } i;$~~

Example →

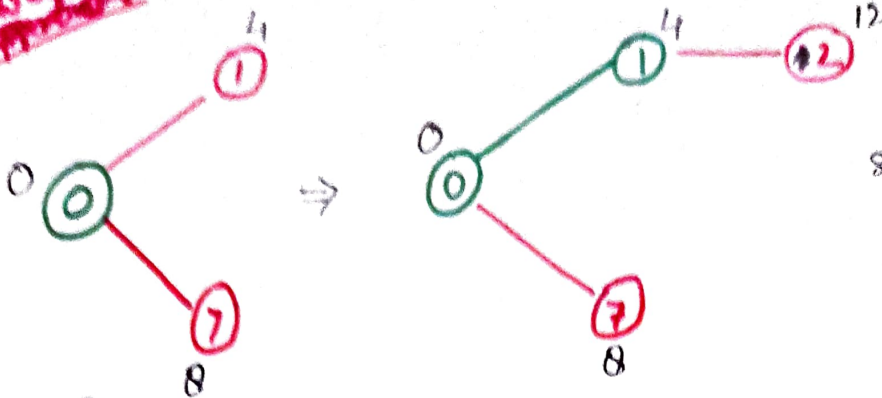


Solution → Dijkstra Algo.



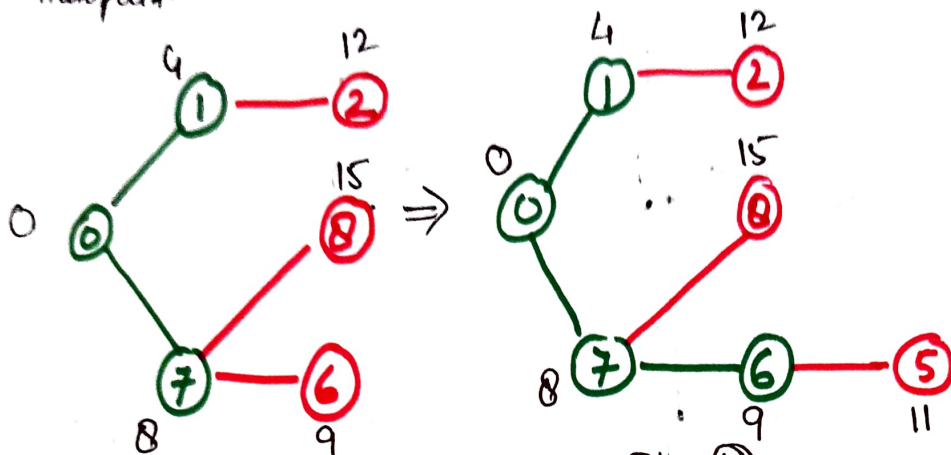
Way Approach

5



step (2)

step (1)
min path.



step (4)

step (3)

