## Lecture of Module 4

Sequential Circuits (Latch/Flip Flop)

## Overview

- Introduction
- Latch
- Flip Flops
- Triggering of Flip Flop
- SR, D, JK, T Flip Flop
- Master Slave Flip Flop
- Characteristic Table and Characteristic Equation
- Excitation Table


## Sequential Logic Circuits

The output state of a "sequential logic circuit" is a function of the following three states, the "present input", the "past input" and/or the "past output". Sequential Logic circuits remember these conditions and stay fixed in their current state until the next clock signal changes.

Sequential logic circuits are generally termed as two state or Bistable devices. Outputs set in one of two basic states, a logic level " 1 " or a logic level " 0 " and will remain "latched" (hence the name is latch) indefinitely in this current state or condition until some other input trigger pulse or signal is applied which will cause the bistable to change its state once again.


## Sequential Logic: Concept

- Sequential Logic circuits remember past inputs and past circuit state.
- Outputs from the system are "fed back" as new inputs.
- The storage elements are circuits that are capable of storing binary information: Memory.

The basic sequential circuit elements can be divided in two categories:

- Level-sensitive (Latches)
- High-level sensitive
- Low-level sensitive
- Edge-triggered (Flip-flops)
- Rising (positive) edge triggered
- Falling (negative) edge triggered
- Dual-edge triggered


## Clock

Sequential circuits can be Asynchronous or Synchronous.
Asynchronous sequential circuits change their states and output values whenever a change in input values occurs. Circuit output can change at any time (clock less).

Synchronous sequential circuits change their states and output values at fixed points of time.
This type of circuits achieves synchronization by using a timing signal called the clock.


Clock generator: Periodic train of clock pulses

## Memory Devices

Latches: A latch is a memory element whose excitation signals control the state of the device. A latch has two stages set and reset. Set stage sets the output to 1. Reset stage set the output to 0 .

Latches are also called level triggered flip flops, because the change on the outputs will follow the changes of the inputs as long as the Enable input is set.

* This causes synchronization problems.

Solution: use latches to create flip-flops that can respond (update) only on specific times (instead of any time).

Flip-flops: A flip-flop is a memory device that has clock signals control the state of the device.

Flip Flops are Edge triggered that change there outputs only at the transition of the clock signal.

## Latch Vs. Flip Flop

Lutches

| Latches are building blocks of sequential |
| :--- |
| circuits and these can be built from logic |
| gates |


| Latch continuously checks its inputs and |
| :--- |
| changes its output correspondingly. |

The latch is sensitive to the duration of the
pulse and can send or receive the data
when the switch is on
It is based on the enable function input
It is a level triggered, it means that the
output of the present state and input of the
next state depends on the level that is
binary input or o.
Flip Flops

| Flip flops are also building blocks oof |
| :--- |
| sequential circuits. But, these can be built |
| from the latches. |
| Flip flop continuously checks its inputs |
| and changes its output correspondingly |
| only at times determined by clocking |
| signal |


| Flipflop is sensitive to a signal change. |
| :--- |
| They can transfer data only at the single |
| next signal change. Flip flops areused as a |
| register. |

It works on the basis of clockpulses
It is an edge triggered, it means that the
output and the nextstate inputchanges
when there is a change inclockpulse
whether it may a +ve or -veclock pulse.

## Synchronous Sequential Circuits: Flip flops as state memory


(a) Block diagram

(b) Timing diagram of clock pulses

- The flip-flops receive their inputs from the combinational circuit and also from a clock signal with pulses that occur at fixed intervals of time, as shown in the timing diagram.


## S-R Latch(NOR version)



Recall...


Time | $R$ | $S$ | $Q$ | $\bar{Q}$ | Comment |
| :--- | :--- | :--- | :--- | :--- |
|  | 0 | $?$ | $?$ | Stored state unknown |
| 0 | 1 | 1 | 0 | "Set" Q to 1 |
| 0 | 0 | 1 | 0 | Now Q "remembers" 1 |
| 1 | 0 | 0 | 1 | "Reset" Q to 0 |
| 0 | 0 | 0 | 1 | Now Q "remembers" 0 |
| 1 | 1 | 0 | 0 | Both go low |
| 0 | 0 | $?$ | $?$ | Unstable! |

## S-R Latch(NAND version)



## S-R Latch(NAND version)



| $x$ | $y$ | NAND |
| :---: | :---: | :---: |
| 0 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

Characteristic Table

## S-R Latch(NAND version)



| $x$ | $y$ | NAND |
| :---: | :---: | :---: |
| 0 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |



Characteristic Table

## S-R Latch Flop(NAND version)



## S-R Latch(NAND version)



| $S^{\prime}$ | $R^{\prime}$ | $Q$ | $Q^{\prime}$ |  |
| :---: | :---: | :---: | :---: | :--- |
| 0 | 0 | 1 | 1 | Forbidden |
| 0 | 1 | 1 | 0 | Set |
| 1 | 0 | 0 | 1 | Reset |
| 1 | 1 | 1 | 0 | Last State |
|  |  | 0 | 1 | Last State |


| $x$ y | NAND | Characteristic Table |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 00 | 1 | $s$ | R | Next state or |
| 01 | 1 | - | o | No chang |
| 10 | 1 |  | 1 | Q $=0$; Res |
| 11 | 0 | 1 | 1 | Undefined |

## S-R Latch(NAND version)



| $S$ | $R$ | Next state of $Q$ |
| :--- | :--- | :--- |
| 0 | 0 | No change |
| 0 | 1 | $Q=0 ;$ Reset state |
| 1 | 0 | $Q=1 ;$ Set state |
| 1 | 1 | Undefined |

Function table
Logic diagram

## S-R Flip Flop with Clock signal


(a) Logic diagram

| $C$ | $S$ | $R$ | Next state of $Q$ |
| :--- | :--- | :--- | :--- |
| 0 | $\times$ | $\times$ | No change |
| 1 | 0 | 0 | No change |
| 1 | 0 | 1 | $Q=0 ;$ Reset state |
| 1 | 1 | 0 | $Q=1 ;$ Set state |
| 1 | 1 | 1 | Undefined |

(b) Function table

Latch is sensitive to input changes ONLY when $\mathrm{C}=1$

## S-R Flip Flop with Clock signal



Logic diagram

| $C$ | $S$ | $R$ | Next state of $Q$ |
| :--- | :--- | :--- | :--- |
| 0 | $\times$ | $\times$ | No change |
| 1 | 0 | 0 | No change |
| 1 | 0 | 1 | $Q=0 ;$ Reset state |
| 1 | 1 | 0 | $Q=1 ;$ Set state |
| 1 | 1 | 1 | Undefined |

Function table

Latch is sensitive to input changes ONLY when $\mathrm{C}=1$

## Characteristic Equation of S-R Flip Flop



| $R$ | $S$ | $Q$ | $\bar{Q}$ |  |
| :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | $Q_{0}$ | $\bar{Q}_{0}$ |  | | Q and $\bar{Q}_{0}$ |
| :--- |
| 0 | 1 | 1 | 0 | are initial conditions |  |
| :--- | :--- | :--- | :--- |
| 1 | 0 | 0 | 1 |$\quad$ (not defined)


| $Q$ | $S$ | $R$ | $Q(t+1)$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | Negative Status |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | Negative Status |

Characteristic Table


Characteristic Equation

## Triggering of Flip Flop



Triggers on this edge
$Q_{0}$ and $\bar{Q}_{0}$
are initial conditions
(not defined)
of the clock pulse


Positive Edge Triggering

Triggers on this edge of the clock pulse


Negative Edge Triggering

## D Flip Flop

- One way to eliminate the undesirable indeterminate state in the RS flip flop is to ensure that inputs S and R are never 1 simultaneously.


| C | D | Next state of Q |
| :--- | :--- | :--- |
| 0 | $X$ | No change |
| 1 | 0 | $\mathrm{Q}=0$; Reset state |
| 1 | 1 | $\mathrm{Q}=1$; Set state |



## Characteristic Equation of D Flip Flop

Characteristics table

| Qn | D | $\mathrm{Q}(\mathrm{n}+1)$ |
| :--- | :--- | :--- |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |


$\mathrm{Q}(\mathrm{n}+1)=\mathrm{D}$

Characteristic Equation

## J-K Flip Flop

- In SR Flip Flop $\mathrm{S}=\mathrm{R}=1$ should be avoided.
$\rightarrow$ To overcome that JK Flip Flop developed.
$\rightarrow$ Both the S and the R inputs of the previous SR bistable have now been replaced by two inputs called the J and K inputs respectively after its inventor Jack Kilby. Then this may equates to: $\mathrm{J}=\mathrm{S}$ and $\mathrm{K}=\mathrm{R}$.
- When $\mathrm{J}=0, \mathrm{~K}=0$, no change in state.
- When $\mathrm{J}=0, \mathrm{~K}=1$, Q will reset.
- When $\mathrm{J}=1, \mathrm{~K}=0$, Q will set.
- When $\mathrm{J}=1, \mathrm{~K}=1$, Toggle i.e $\mathrm{Q}^{\prime}{ }_{\mathrm{n}}$


Symbol



- When $\mathrm{J}=1, \mathrm{~K}=1$, Toggle i.e $\mathrm{Q}_{\mathrm{n}}$
- For JK flip-flop if J, K and Clock are equal to 1 the state of flip-flop keeps on toggling which leads to uncertainty in determining the output of the flipflop. This problem is called Race around the condition.
- This can be avoided by
* Using Edge triggering of JK Flip Flop
* Enhancing the propagation delay
* Using Master-Slave Flip Flop



## Master-Slave Flip Flop

- Master-slave flip flop is designed using two separate flip flops. Out of these, one acts as the master and the other as a slave.
- The J-K flip flops are presented in a series connection.
- The output of the master J-K flip flop is fed to the input of the slave J-K flip flop.
- The output of the slave J-K flip flop is given as a feedback to the input of the master J-K flip flop.
- The clock pulse [Clk] is given to the master J-K flip flop and it is sent through a NOT Gate and thus inverted before passing it to the slave J-K flip flop.
- It avoids the race around condition of J-K Flip Flop



## Characteristic Equation of J-K Flip Flop

Characteristic Table

| $Q$ | $J$ | $K$ | $Q(t+1)$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 0 |

JK flip-flop


K Map

$$
Q(t+1)=J Q^{\prime}+K^{\prime} Q
$$

Characteristic Equation

## T Flip Flop

- We can construct the "T Flip Flop" by making changes in the "JK Flip Flop".
- The "T Flip Flop" has only one input, which is constructed by connecting the input of JK Flip Flop.
- This single input is called T.
- Sometimes the "T Flip Flop" is referred to as single input "JK Flip Flop".
- In T flip flop, "T" defines the term "Toggle"

| $Q$ | $T$ | $Q(T+1)$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| $\mathbf{0}$ | $\mathbf{1}$ | 1 |
| $\mathbf{1}$ | $\mathbf{0}$ | 1 |
| $\mathbf{1}$ | $\mathbf{1}$ | 0 |



## Characteristic Equation of T Flip Flop

## Characteristic Table

| $Q$ | $T$ | $Q(t+1)$ |
| :--- | :--- | :--- |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |
| $T$ flip-flop |  |  |


$Q(+1)=T Q^{\prime}+T^{\prime} Q$
Characteristic Equation

## Excitation Table

- The characteristic table is useful for analysis and for defining the operation of flip flop.
- It specifies the next state when the inputs and present state are known.
- During design process we usually know the transition from present state to next state.
- So, we want to know the flip flop input conditions that will cause the required transition.
- Therefore, we need a table that lists the required inputs for a given change of states.
- Such table is called as Excitation Table.

| JK flip-flop |  |  |  |
| :---: | :---: | :---: | :---: |
| $Q(t)$ | $Q(t+1)$ | J | K |
| 0 | 0 | 0 | x |
| 0 | 1 | 1 | x |
| 1 | 0 | x | 1 |
| 1 | 1 | x | 0 |


| T flip-flop |  |  |  |
| :---: | :---: | :---: | :---: |
| $Q(t)$ | $Q(t+1)$ | $T$ |  |
| 0 | 0 | 0 |  |
| 0 | 1 | 1 |  |
| 1 | 0 | 1 |  |
| 1 | 1 | 0 |  |


| SR Flip-flop |  |  |  |
| :---: | :---: | :---: | :---: |
| Q(t) | Q(t+1) | S | R |
| 0 | 0 | 0 | X |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 1 | X | 0 |


| DFlip-flop |  |  |  |
| :---: | :---: | :---: | :---: |
| $Q(t)$ | $Q(t+1)$ | $D$ |  |
| 0 | 0 | 0 |  |
| 0 | 1 | 1 |  |
| 1 | 0 | 0 |  |
| 1 | 1 | 1 |  |

## Sequential Circuit Design

Example:

| Count Sequences |  |  | Flip Flop Inputs |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| A | B | C | TA | TB | TC |
| 0 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 1 | 0 | 1 | 1 |
| 0 | 1 | 0 | 0 | 0 | 1 |
| 0 | 1 | 1 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 | 1 |



## Example:

Let the state equations are:
$\mathrm{A}(\mathrm{t}+1)=\mathrm{A}^{\prime} \mathrm{B}^{\prime} \mathrm{CD}+\mathrm{A}^{\prime} \mathrm{B}^{\prime} \mathrm{C}+\mathrm{ACD}+\mathrm{AC}^{\prime} \mathrm{D}^{\prime}$
$B(t+1)=A^{\prime} C+C D^{\prime}+A^{\prime} B^{\prime}$
$C(t+1)=B$
$\mathrm{D}(\mathrm{t}+1)=\mathrm{D}^{\prime}$
The above equation can be rearranged in the form of characteristic equation of $\mathrm{J}-\mathrm{K}$ flip flop.
Characteristic equation of J-K flip flop is

$$
Q(t+1)=J Q^{\prime}+K^{\prime} Q
$$

$\mathrm{A}(\mathrm{t}+1)=\mathrm{A}^{\prime} \mathrm{B}^{\prime} \mathrm{CD}+\mathrm{A}^{\prime} \mathrm{B}^{\prime} \mathrm{C}+\mathrm{ACD}+\mathrm{AC}^{\prime} \mathrm{D}^{\prime}$

$$
=\left(\mathrm{B}^{\prime} \mathrm{CD}+\mathrm{B}^{\prime} \mathrm{C}\right) \mathrm{A}^{\prime}+\left(\mathrm{CD}+\mathrm{C}^{\prime} \mathrm{D}^{\prime}\right) \mathrm{A}
$$

So, $J=\mathrm{B}^{\prime} \mathrm{CD}+\mathrm{B}^{\prime} \mathrm{C}=\mathrm{B}^{\prime} \mathrm{C}$

$$
\mathrm{K}=\left(\mathrm{CD}+\mathrm{C}^{\prime} \mathrm{D}^{\prime}\right)^{\prime}=\mathrm{C}^{\prime} \mathrm{D}+\mathrm{CD}^{\prime}
$$

$\mathrm{B}(\mathrm{t}+1)=\mathrm{A}^{\prime} \mathrm{C}+\mathrm{CD}^{\prime}+\mathrm{A}^{\prime} \mathrm{BC}^{\prime}$

$$
\begin{aligned}
& =\left(\mathrm{A}^{\prime} \mathrm{C}+\mathrm{CD}^{\prime}\right)\left(\mathrm{B}+\mathrm{B}^{\prime}\right)+\mathrm{A}^{\prime} \mathrm{BC} \mathrm{C}^{\prime} \\
& =\left(\mathrm{A}^{\prime} \mathrm{C}+\mathrm{CD}^{\prime}\right) \mathrm{B}^{\prime}+\left(\mathrm{A}^{\prime} \mathrm{C}+\mathrm{CD}^{\prime}\right) \mathrm{B}+\mathrm{A}^{\prime} \mathrm{BC} \mathrm{C}^{\prime} \\
& =\left(\mathrm{A}^{\prime} \mathrm{C}+\mathrm{CD}^{\prime}\right) \mathrm{B}^{\prime}+\left(\mathrm{A}^{\prime} \mathrm{C}+\mathrm{CD}^{\prime}+\mathrm{A}^{\prime} \mathrm{C}^{\prime}\right) \mathrm{B}
\end{aligned}
$$

So, $J=A^{\prime} C+C D^{\prime}$

$$
\mathrm{K}=\left(\mathrm{A}^{\prime} \mathrm{C}+\mathrm{CD}^{\prime}+\mathrm{A}^{\prime} \mathrm{C}^{\prime}\right)^{\prime}=\mathrm{AC}^{\prime}+\mathrm{AD}
$$

$$
\begin{aligned}
& \mathrm{C}(\mathrm{t}+1)=\mathrm{B}=\mathrm{B}\left(\mathrm{C}+\mathrm{C}^{\prime}\right)=\mathrm{BC}^{\prime}+\mathrm{BC} \\
& \text { So, } \mathrm{J}=\mathrm{B} \\
& \quad \mathrm{~K}=\mathrm{B}^{\prime} \\
& \mathrm{D}(\mathrm{t}+1)=\mathrm{D}^{\prime}=(1) \mathrm{D}^{\prime}+(0) \mathrm{D} \\
& \text { So, } J=1 \\
& \quad \mathrm{~K}=1
\end{aligned}
$$

So, finally

$$
\begin{array}{ll}
\mathrm{J} A=\mathrm{B}^{\prime} \mathrm{C} & \mathrm{KA}=\mathrm{C}^{\prime} \mathrm{D}+\mathrm{CD}^{\prime} \\
\mathrm{JB}=\mathrm{A}^{\prime} \mathrm{C}+\mathrm{CD}^{\prime} & \mathrm{KB}=\mathrm{AC}^{\prime}+\mathrm{AD} \\
\mathrm{JC}=\mathrm{B} & \mathrm{KC}=\mathrm{B}^{\prime} \\
\mathrm{J} D=1 & \mathrm{KD}=1
\end{array}
$$

## Setup time and Hold time

The Clocking event can be either from low to high or from high to low. The input signal around clocking event must remain unchanged in order to have a correct effect on the outcome of the new state.
$\rightarrow T_{\mathrm{s}}$ : the minimum time interval preceding the clocking event the input signal must remain available and unchanged.
$\rightarrow T_{\mathrm{h}}$ : the minimum time interval after edge of the clocking event, the input signals must remain
 unchanged

## Applications

Flip flops will find their use in many of the fields in digital electronics. Flip flops are the main components of sequential circuits. Particularly, edge triggered flip flops are very resourceful devices that can be used in wide range of applications like storing of binary data and transferring binary data from one location to other etc. Some of the most common applications of flip flops are

- Shift Register
- Counter
- Storage Registers
- Memory
- Data transfer
- Frequency Dividers etc.

Sequential Circuits (Shift Register)

## Overview

- Register
- Shift Register
- Types of Shift Register
- Bidirectional Shift Register
- Universal Shift Register
- Typical ICs for Shift register


## Register

- The filp flops are essential component in clocked sequential circuits.
- Circuits that include filp flops are usually classified by the function they perform.
- Two such circuits are registers and counters.
- An n -bit register consists of a group of n flip flops capable of storing n bits of binary information.
- So, Register is a collection of flip flops.
- A flip flop is used to store single bit digital data. For storing a large number of bits, the storage capacity is increased by grouping more than one flip flops.
- It is used to perform simple data storage, movement, manipulation and processing operations (e.g. load, increment, shift, add, etc.)
- The computer processes data by performing operations on registers, e.g. ADD A, B where A and B are the registers.
- A register capable of shifting its binary information in one or both direction is called a shift register.
- All flip-flops receive common clock pulses, which activate the shift from one stage to the next.


## Shift Register

- The simplest possible shift register is one that uses only flip-flops, as shown in Fig.
- Each clock pulse shifts the contents of the register one bit position to the right.
- The serial input determines what goes into the leftmost flip-flop during the shift.
- The serial output is taken from the output of the rightmost flip-flop.
- Following are the four types of shift registers based on applying inputs and accessing of outputs.
- Serial In - Serial Out shift register
- Serial In - Parallel Out shift register
- Parallel In - Serial Out shift register
- Parallel In - Parallel Out shift register



## Serial In - Serial Out (SISO) shift register

- The shift register, which allows serial input and produces serial output is known as Serial In - Serial Out (SISO) shift register.
- This block diagram consists of three D flip-flops, which are cascaded. That means, output of one D flip-flop is connected as the input of next D flip-flop.
- All these flip-flops are synchronous with each other since, the same clock signal is applied to each one.
- In this shift register, we can send the bits serially from the input of left most D flip-flop. Hence, this input is also called as serial input.
- For every positive edge triggering of clock signal, the data shifts from one stage to the next. So, we can receive the bits serially from the output of right most D flip-flop. Hence, this output is also called as serial output.



## Serial In - Parallel Out (SIPO) shift register

- The shift register, which allows serial input and produces parallel output is known as Serial In Parallel Out (SIPO) shift register.
- In this shift register, we can send the bits serially from the input of left most D flip-flop. Hence, this input is also called as serial input.
- For every positive edge triggering of clock signal, the data shifts from one stage to the next.
- In this case, we can access the outputs of each D flip-flop in parallel. So, we will get parallel outputs from this shift register.



## Parallel In - Serial Out (PISO) shift register

- In the 'Parallel In Serial Out'" register, the data is entered in a parallel way, and the outcome comes serially.
- The shift mode and the load mode are the two modes in which the 'PISO" circuit works.



## Parallel In - Parallel Out (PIPO) shift register

- In "Parallel In Parallel Out", the inputs and the outputs come in a parallel way in the register.
$\rightarrow$ The inputs $\mathrm{B}_{0}, \mathrm{~B}_{1}, \mathrm{~B}_{2}$, and $\mathrm{B}_{3}$, are directly passed to the data inputs $\mathrm{D}_{0}, \mathrm{D}_{1}, \mathrm{D}_{2}$, and $\mathrm{D}_{3}$ of the respective flip flop.
- The bits of the binary input is loaded to the flip flops when the negative clock edge is applied. The clock pulse is required for loading all the bits. At the output side, the loaded bits appear.



## Bidirectional Shift Register

- Below is the diagram of 4-bit "bidirectional" shift register where $\mathbf{D}_{\mathbf{R}}$ is the "serial right shift data input", $\mathbf{D}_{\mathrm{L}}$ is the 'left shift data input", and M is the "mode select input".



## Universal Shift Register

- A shift-right control to enable the shift operation and the serial input and output lines associated with the shift right.
- A shift-left control to enable the shift operation and the serial input and output lines associated with the shift left.
- A parallel-load control to enable a parallel transfer and the $\boldsymbol{n}$ input lines associated with the parallel transfer.
- If the Shift register has the capability of

Serial In - Serial Out
Serial In - Parallel Out
Parallel In - Serial Out
Parallel In - Parallel Out
and act as Bidirectional shift register is referred as a universal shift register.

- Shift registers are often used to interface digital system situated remotely from each other. If the distance is far, it will be expensive to use $n$ lines to transmit the $n$ bits in parallel.
- Transmitter performs a parallel-to-serial conversion of data and the receiver does a serial-to-parallel conversion.



## Universal Shift Register using MUX



| Mode Control |  |  |
| :--- | :--- | :--- |
|  |  |  |
| $S_{1}$ | $S_{0}$ | Register Operation |
| 0 | 0 | No Change |
| 0 | 1 | Shitright |
| 1 | 0 | Shititleft |
| 1 | 1 | Parallel load |

## Typical ICs for Shift register

- Commonly available SISO IC is 74HC595, which is 8-bit.
- Commonly available SIPO IC's include the standard 8-bit 74LS164, 74LS594.
- Commonly available PISO IC is 74 HC 166 , which is 8 -bit.
- Commonly available PIPO IC's include the standard 8-bit M54HC195, M74HC195.
- Today, there are many high speed bi-directional or "universal" type Shift Registers available such as the TTL 74LS194, 74LS195 or the CMOS 4035 which are available as 4-bit multi-function devices that can be used in either serial-to-serial, left shifting, right shifting, serial-to-parallel, parallel-to-serial, or as a parallel-to-parallel multifunction data register, hence their name "Universal".
- The 74HC299 is an 8-bit Universal Shift register.
- The 74S299 is an 8-bit Universal Shift and Storage Register.

